

# Domain decomposition methods for compressed sensing

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## Abstract

We present several domain decomposition algorithms for sequential and parallel minimization of functionals formed by a discrepancy term with respect to data and total variation constraints. The convergence properties of the algorithms are analyzed. We provide several numerical experiments, showing the successful application of the algorithms for the restoration 1D and 2D signals in interpolation/inpainting problems respectively, and in a compressed sensing problem, for recovering piecewise constant medical-type images from partial Fourier ensembles.

## Introduction

In concrete applications for image processing, one might be interested to recover at best a digital image provided only partial linear or nonlinear measurements, possibly corrupted by noise. Given the observation that natural and man-made images are characterized by a relatively small number of edges and extensive relatively uniform parts, one may want to help the reconstruction by imposing that the interesting solution is the one which matches the given data and has also a few discontinuities localized on sets of lower dimension.

In the context of compressed sensing [2, 8], it has been clarified the fundamental role of minimizing  $\ell_1$ -norms in order to promote sparse solutions. This understanding furnishes an important interpretation of total variation minimization, i.e., the minimization of the  $\ell_1$ -norm of derivatives [13], as a regularization technique for image restoration. Several numerical strategies to perform efficiently total variation minimization have been proposed in the literature. We list a few of the relevant ones, ordered by their chronological appearance:

(i) the approach of Chambolle and Lions [4] by re-weighted least squares, see also [6] for generalizations and refinements in the context of compressed sensing;

(ii) variational approximation via local quadratic functionals as in the work of Vese et al. [14, 1];

(iii) iterative thresholding algorithms based on projections onto convex sets as in the work of Chambolle [3] as well as in the work of Combettes-Wajs [5] and Daubechies et al. [7];

(iv) iterative minimization of the Bregman distance as in the work of Osher et al. [12];

(v) the approach proposed by Nesterov [11] and its modifications by Weiss et al. [15].

These approaches differ significantly, and it seems that the ones collected in the groups iv) and v) do show presently the best performances in practice. However, none of the mentioned methods is able to address in real-time, or at least in an acceptable computational time, extremely large problems, such as 4D imaging (spatial plus temporal dimensions) from functional magnetic-resonance in nuclear medical imaging, astronomical imaging or global terrestrial seismic tomography. For such large scale simulations we need to address methods which allow us to reduce the problem to a finite sequence of sub-problems of more manageable size, perhaps by one of the methods listed above. With this aim we introduced subspace correction and domain decomposition methods both for  $\ell_1$ -norm and total variation minimizations [9, 10]. Due to the nonadditivity of the total variation with respect to a domain decomposition (the total variation of a function on the whole domain equals the sum of the total variations on the subdomains plus the size of the jumps at the interfaces), one encounters additional difficulties in showing convergence of such decomposition strategies to global minimizers.

In this paper we review concisely both nonoverlapping and overlapping domain decomposition methods for total variation minimization and we provide their properties of convergence to global minimizers. Moreover, we show numerical applications in classical problems of signal and image processing, such as signal interpolation and image inpainting. We further include applications in the context of compressed sensing for recovering piecewise constant medical-type images from partial Fourier ensembles [2].

## 2 Notations and preliminaries

Since we are interested to a discrete setting, we define the domain of our multivariate digital signal  $\Omega = \{x_1^1 < \dots < x_{N_1}^1\} \times \dots \times \{x_1^d < \dots < x_{N_d}^d\} \subset \mathbb{R}^d$ ,  $d \in \mathbb{N}$  and we consider the signal space  $\mathcal{H} = \mathbb{R}^{N_1 \times N_2 \times \dots \times N_d}$ , where  $N_i \in \mathbb{N}$

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for  $i = 1, \dots, d$ . For  $u \in \mathcal{H}$  we write  $u = u(x_i)_{i \in \mathcal{I}}$  with index set  $\mathcal{I} := \prod_{k=1}^d \{1, \dots, N_k\}$  and  $u(x_i) = u(x_{i_1}^1, \dots, x_{i_d}^d)$  where  $i_k \in \{1, \dots, N_k\}$  and  $x_i \in \Omega$ . Then, for  $g \in \mathbb{R}^N$ , the  $\ell_p$ -norm is given by  $\|g\|_p = \left( \sum_{k=1}^N |g_k|^p \right)^{1/p}$ ,  $1 \leq p < \infty$ , and for  $u \in \mathcal{H}$  the discrete gradient  $\nabla u$  is given by  $(\nabla u)(x_i) = ((\nabla u)^1(x_i), \dots, (\nabla u)^d(x_i))$  with  $(\nabla u)^j(x_i) = u(x_{i_1}^1, \dots, x_{i_{j+1}}^j, \dots, x_{i_d}^d) - u(x_{i_1}^1, \dots, x_{i_j}^j, \dots, x_{i_d}^d)$  if  $i_j < N_j$ , and  $(\nabla u)^j(x_i) = 0$  if  $i_j = N_j$ , for all  $j = 1, \dots, d$  and for all  $\mathbf{i} = (i_1, \dots, i_d) \in \mathcal{I}$ . The total variation of  $u \in \mathcal{H}$  in the discrete setting is then defined as  $|D(u)|(\Omega) = \sum_{i \in \mathcal{I}} |(\nabla u)(x_i)|$ ,

with  $|y| = \sqrt{y_1^2 + \dots + y_d^2}$  for every  $y = (y_1, \dots, y_d) \in \mathbb{R}^d$ . We define the scalar product of  $u, v \in \mathcal{H}$  as usual,  $\langle u, v \rangle_{\mathcal{H}} = \sum_{i \in \mathcal{I}} u(x_i)v(x_i)$ , and the scalar product of  $p, q \in \mathcal{H}^d$  as  $\langle p, q \rangle_{\mathcal{H}^d} = \sum_{i \in \mathcal{I}} p^1(x_i)q^1(x_i) + \dots + p^d(x_i)q^d(x_i)$ . Further we introduce a discrete divergence  $\operatorname{div} : \mathcal{H}^d \rightarrow \mathcal{H}$  defined, by analogy with the continuous setting, by  $\operatorname{div} = -\nabla^*$  ( $\nabla^*$  is the adjoint of the gradient  $\nabla$ ). In the following we denote with  $\pi_K$  the orthogonal projection onto a closed convex set  $K$ .

## 2.1 Projections onto convex sets

With these notation, we define the closed convex set

$$K := \{ \operatorname{div} p : p \in \mathcal{H}^d, |p_i| \leq 1 \text{ for all } i \in \mathcal{I} \}.$$

We briefly recall here an algorithm proposed by Chambolle in [3] in order to compute the projection onto  $\alpha K$ . The following semi-implicit gradient descent algorithm is given to approximate  $\pi_{\alpha K}(g)$ : Choose  $\tau > 0$ , let  $p^{(0)} = 0$  and, for any  $n \geq 0$ , iterate

$$p_i^{(n+1)} = \frac{p_i^{(n)} + \tau(\nabla(\operatorname{div} p^{(n)} - g/\alpha))_i}{1 + \tau |(\nabla(\operatorname{div} p^{(n)} - g/\alpha))_i|}. \quad (1)$$

For  $\tau > 0$  sufficiently small, the iteration  $\alpha \operatorname{div} p^{(n)}$  converges to  $\pi_{\alpha K}(g)$  as  $n \rightarrow \infty$  (compare [3, Theorem 3.1]).

## 2.2 Setting of the problem

Given a model linear operator  $T : \mathcal{H} \rightarrow \mathbb{R}^K$ , we are considering the following discrete minimization problem

$$\arg \min_{u \in \mathcal{H}} \{ \mathcal{J}(u) := \|Tu - g\|_2^2 + 2\alpha |Du|(\Omega) \} \quad (2)$$

where  $g \in \mathbb{R}^K$  is a given datum and  $\alpha > 0$  is a fixed regularization parameter. Note that, up to rescaling the parameter  $\alpha$  and the datum  $g$ , we can always assume  $\|T\| < 1$ . Moreover, in order to ensure existence of solutions, we assume  $1 \notin \ker(T)$ . For both nonoverlapping and overlapping domain decompositions, we will consider a linear sum  $\mathcal{H} = V_1 + V_2$  with respect to two subspaces  $V_1, V_2$  defined by a suitable decomposition of the physical domain  $\Omega$ . We restrict our discussion to two subspaces, but the analysis can be extended in a straightforward way to multiple subspaces. With this splitting we want to minimize  $\mathcal{J}$  by suitable instances of the following alternating algorithm: Pick an initial

$V_1 + V_2 \ni u_1^{(0)} + u_2^{(0)} := u^{(0)} \in \mathcal{H}$ , for example  $u^{(0)} = 0$ , and iterate

$$\begin{cases} u_1^{(n+1)} \approx \arg \min_{v_1 \in V_1} \mathcal{J}(v_1 + u_2^{(n)}) \\ u_2^{(n+1)} \approx \arg \min_{v_2 \in V_2} \mathcal{J}(u_1^{(n+1)} + v_2) \\ u^{(n+1)} := u_1^{(n+1)} + u_2^{(n+1)}. \end{cases}$$

## 3 Nonoverlapping domain decomposition methods

Let us consider the disjoint domain decomposition  $\Omega = \Omega_1 \cup \Omega_2$  and  $\Omega_1 \cap \Omega_2 = \emptyset$  and the corresponding spaces  $V_j = \{u \in \mathcal{H} : \operatorname{supp}(u) \subset \Omega_j\}$ , for  $j = 1, 2$ . Note that  $\mathcal{H} = V_1 \oplus V_2$ . It is useful to us to introduce an auxiliary functional  $\mathcal{J}_1^s$ , called the *surrogate functional* of  $\mathcal{J}$ : For  $j \in \{1, 2\}$  and  $\check{j} \in \{1, 2\} \setminus \{j\}$ , assume  $a, u_j \in V_j$  and  $u_{\check{j}} \in V_{\check{j}}$  and define

$$\mathcal{J}_j^s(u_j + u_{\check{j}}, a) := \mathcal{J}(u_j + u_{\check{j}}) + \|u_j - a\|_{\mathcal{H}}^2 - \|T(u_j - a)\|_{\mathbb{R}^K}^2. \quad (3)$$

As it will be clarified later, the minimization of  $\mathcal{J}_j^s(u_j + u_{\check{j}}, a)$  with respect to  $u_j$  and for fixed  $u_{\check{j}}, a$  is an operation which can be realized more easily than the direct minimization of the parent functional  $\mathcal{J}(u_j + u_{\check{j}})$  for the sole  $u_{\check{j}}$  fixed.

### 3.1 Sequential algorithm

In the following we denote  $u_j = \pi_{V_j} u$  the orthogonal projection onto  $V_j$ , for  $j = 1, 2$ . Let us explicitly express the algorithm as follows: Pick an initial  $V_1 \oplus V_2 \ni u_1^{(0,L)} + u_2^{(0,M)} := u^{(0)} \in \mathcal{H}$ , for example  $u^{(0)} = 0$ , and iterate

$$\begin{cases} \begin{cases} u_1^{(n+1,0)} = u_1^{(n,L)} \text{ and for } \ell = 0, \dots, L-1 \\ u_1^{(n+1,\ell+1)} = \arg \min_{u_1 \in V_1} \mathcal{J}_1^s(u_1 + u_2^{(n,M)}, u_1^{(n+1,\ell)}) \end{cases} \\ \begin{cases} u_2^{(n+1,0)} = u_2^{(n,M)} \text{ and for } m = 0, \dots, M-1 \\ u_2^{(n+1,m+1)} = \arg \min_{u_2 \in V_2} \mathcal{J}_2^s(u_1^{(n+1,L)} + u_2, u_2^{(n+1,m)}) \end{cases} \\ u^{(n+1)} := u_1^{(n+1,L)} + u_2^{(n+1,M)}. \end{cases} \quad (4)$$

Note that we do prescribe a finite number  $L$  and  $M$  of inner iterations for each subspace respectively.

### 3.2 Parallel algorithm

The parallel version of the previous algorithm reads as follows: Pick an initial  $V_1 \oplus V_2 \ni u_1^{(0,L)} + u_2^{(0,M)} := u^{(0)} \in \mathcal{H}$ , for example  $u^{(0)} = 0$ , and iterate

$$\begin{cases} \begin{cases} u_1^{(n+1,0)} = u_1^{(n,L)} \text{ and for } \ell = 0, \dots, L-1 \\ u_1^{(n+1,\ell+1)} = \arg \min_{u_1 \in V_1} \mathcal{J}_1^s(u_1 + u_2^{(n,M)}, u_1^{(n+1,\ell)}) \end{cases} \\ \begin{cases} u_2^{(n+1,0)} = u_2^{(n,M)} \text{ and for } m = 0, \dots, M-1 \\ u_2^{(n+1,m+1)} = \arg \min_{u_2 \in V_2} \mathcal{J}_2^s(u_1^{(n,L)} + u_2, u_2^{(n+1,m)}) \end{cases} \\ u^{(n+1)} := \frac{u_1^{(n+1,L)} + u_2^{(n+1,M)} + u^{(n)}}{2}. \end{cases} \quad (5)$$

Note that  $u^{(n+1)}$  is the average of the current iteration and the previous one as it is the case for successive overrelaxation methods (SOR) in classical numerical linear algebra.

## 4 Overlapping domain decomposition methods

Let us consider the overlapping domain decomposition  $\Omega = \Omega_1 \cup \Omega_2$  and  $\Omega_1 \cap \Omega_2 \neq \emptyset$  and the corresponding spaces  $V_j = \{u \in \mathcal{H} : \text{supp}(u) \subset \Omega_j\}$ , for  $j = 1, 2$ . Note that now  $\mathcal{H} = V_1 + V_2$  is not anymore a direct sum of  $V_1$  and  $V_2$ , but just a linear sum of subspaces. We define the internal boundaries  $\Gamma_j = \partial\Omega_j \cap \Omega_{\check{j}}$ ,  $j \in \{1, 2\}$  and  $\check{j} \in \{1, 2\} \setminus \{j\}$  (see Figure 1).

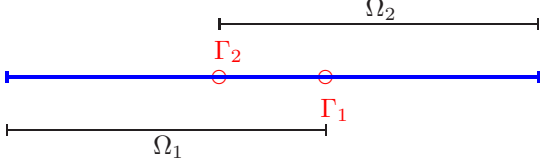


Figure 1: Overlapping domain decomposition and internal boundaries.

### 4.1 Sequential algorithm

Associated to the decomposition  $\{\Omega_j : j = 1, 2\}$  let us fix a partition of unity  $\{\chi_j : j = 1, 2\}$ , i.e.,  $\chi_1 + \chi_2 = 1$ ,  $\text{supp}(\chi_j) \subset \Omega_j$ , and  $\chi_j|_{\Gamma_j} = 0$ . Let us explicitly express the algorithm now as follows: Pick an initial  $V_1 + V_2 \ni \tilde{u}_1^{(0)} + \tilde{u}_2^{(0)} := u^{(0)} \in \mathcal{H}$ , for example  $u^{(0)} = 0$ , and iterate

$$\begin{cases} \begin{cases} u_1^{(n+1,0)} = \tilde{u}_1^{(n)} \text{ and for } \ell = 0, \dots, L-1 \\ u_1^{(n+1,\ell+1)} = \arg \min_{\substack{u_1 \in V_1 \\ u_1|_{\Gamma_1}=0}} \mathcal{J}_1^s(u_1 + \tilde{u}_2^{(n)}, u_1^{(n+1,\ell)}) \end{cases} \\ \begin{cases} u_2^{(n+1,0)} = \tilde{u}_2^{(n)} \text{ and for } m = 0, \dots, M-1 \\ u_2^{(n+1,m+1)} = \arg \min_{\substack{u_2 \in V_2 \\ u_2|_{\Gamma_2}=0}} \mathcal{J}_2^s(u_1^{(n+1,L)} + u_2, u_2^{(n+1,m)}) \end{cases} \\ u^{(n+1)} := u_1^{(n+1,L)} + u_2^{(n+1,M)} \\ \tilde{u}_1^{(n+1)} := \chi_1 \cdot u^{(n+1)} \\ \tilde{u}_2^{(n+1)} := \chi_2 \cdot u^{(n+1)} \end{cases} \quad (6)$$

A few technical tricks are additionally required for the boundedness of the iterations in algorithm (6) with respect to the nonoverlapping version. First of all the local minimizations are restricted to functions which vanish on the internal boundaries. Moreover since  $u^{(n)}$  is formed as a sum of local components  $u_1^{(n)}, u_2^{(n)}$  which are not uniquely determined on the overlapping part, we introduced a suitable correction by means of the partition of unity  $\{\chi_j : j = 1, 2\}$  in order to enforce the uniform boundedness of the sequences of the local iterations  $u_1^{(n)}, u_2^{(n)}$ . With similar minor modifications, we can analogously formulate a parallel version of this algorithm as in (5).

### 4.2 The solution of the local iterations

The inner iterations

$$u_j^{(n+1,\ell+1)} = \arg \min_{\substack{u_j \in V_j \\ \Gamma_j u_j = 0}} \mathcal{J}_j^s(u_j + u_j^{(n)}, u_j^{(n+1,\ell)}),$$

where  $\Gamma_j u_j = 0$  is some linear constraint, are crucial for the concrete realizability of the algorithm. (Note that in the case

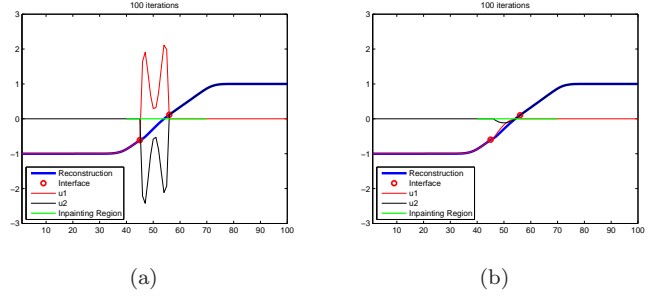


Figure 2: Here we present two numerical experiments related to the interpolation of a 1D signal by total variation minimization, provided only information only out of an interval (indicated in green color in the figures). On the left we show an application of algorithm (6) when no correction with the partition of unity is provided. In this case, the sequence of the local iterations  $u_1^{(n)}, u_2^{(n)}$  is unbounded. On the right we show an application of algorithm (6) with the use of the partition of unity which enforces also the uniform boundedness of the local iterations  $u_1^{(n)}, u_2^{(n)}$ .

of the nonoverlapping decomposition there is no additional linear constraint, whereas for the overlapping case we ask for the trace condition  $\Gamma_j u_j = u_j|_{\Gamma_j} = 0$ .) Such iteration can be explicitly computed

$$u_j^{(n+1,\ell+1)} = [I - \pi_{\alpha K}] \left( u_j^{(n+1,\ell)} + \pi_{V_j} T^* (g - T u_j - T u_j^{(n+1,\ell)} + u_j - \eta^{(n+1,\ell)}) - u_j, \right.$$

for a suitable Lagrange multiplier  $\eta^{(n+1,\ell)}$  which has the role of enforcing the linear constraints  $u_j \in V_j$  and  $\Gamma_j u_j = 0$ ;  $\eta^{(n+1,\ell)}$  can be approximated by an iterative algorithm, see [10, Proposition 4.6] for details. Note that we have to implement repeatedly the projection  $\pi_{\alpha K}$  for which the Chambolle's algorithm (1) is used. More efficient algorithms can also be used such as iterative Bregman distance methods [12] or Nesterov's algorithm [11].

## 5 Convergence properties

These algorithms share common convergence properties, which are listed in the following theorem.

**Theorem.** (Convergence properties) *The algorithms (4), (5), and (6) produce a sequence  $(u^{(n)})_{n \in \mathbb{N}}$  in  $\mathcal{H}$  with the following properties:*

(i)  $\mathcal{J}(u^{(n)}) > \mathcal{J}(u^{(n+1)})$  for all  $n \in \mathbb{N}$  (unless  $u^{(n)} = u^{(n+1)}$ );

(ii)  $\lim_{n \rightarrow \infty} \|u^{(n+1)} - u^{(n)}\|_2 = 0$ ;

(iii) *the sequence  $(u^{(n)})_{n \in \mathbb{N}}$  has subsequences which converge in  $\mathcal{H}$ ; if  $(u^{(n_k)})_{k \in \mathbb{N}}$  is a converging subsequence, and  $u^{(\infty)}$  is its limit, then  $u^{(\infty)}$  is always a minimizer of  $\mathcal{J}$  in the case of algorithm (6) (overlapping case), whereas for algorithms (4), (5) (sequential and parallel nonoverlapping cases) this can be ensured under certain sufficient technical conditions, see [10, Theorem 5.1 and Theorem 6.1] for details.*

